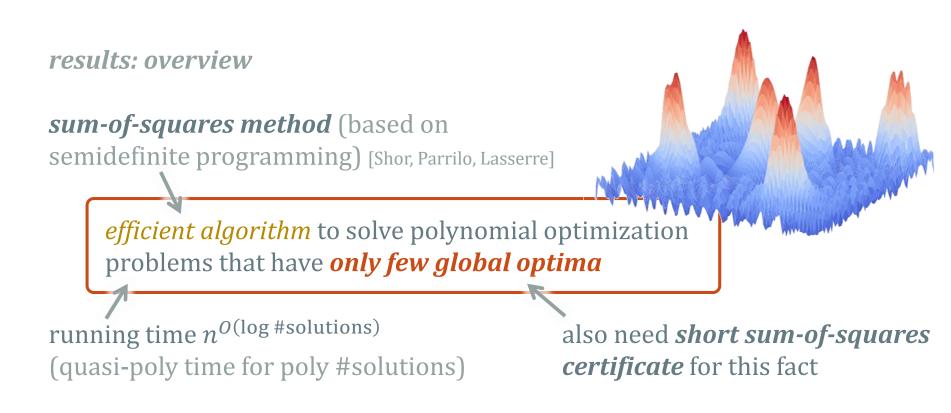
Dictionary learning and tensor decomposition via the sum-of-squares method

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**# bad local optima** can be exponential  $\rightarrow$  local-search algorithms fail

applications: unsupervised learning problems tend to have this property
 identifiability: data uniquely determines parameters of model
 our work: notion of constructive identifiability proofs that
 implies efficient inference algorithms

# results: tensor decomposition

for all constants  $\sigma \ge 1$  and  $\varepsilon > 0$ , exist constants  $d \ge 1$  and  $\tau > 0$ given tensor  $T \in \mathbb{R}^{n^d}$  of the form  $T = \sum_{i=1}^m a_i^{\otimes d} + Z$  with  $||a_i|| = 1$ can recover set  $\approx_{\varepsilon} \{\pm a_1, \dots, \pm a_m\}$  in time  $n^{O(\log n)}$ , whenever  $\|\sum_i a_i a_i^{\mathsf{T}}\|_{\text{spectral}} \le \sigma$  and  $\|Z\|_{\text{spectral}} \le \tau$ 

*comparison to previous algorithms* [Jennrich'70, Bhaskara-Charikar-Moitra-Vijayaraghavan'13, Anandkumar-Ge-Hsu-Kakade'12]

pros:

**tolerate constant spectral error** (*before:* inverse polynomial error) no restrictions on vectors (*before:* incoherence or similar) cons:

running time (*but:* techniques help for faster alg's [Hopkins-Schramm-Shi-S.'15+]) only constant accuracy (*but:* could combine with local search)

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connection to polynomial optimization

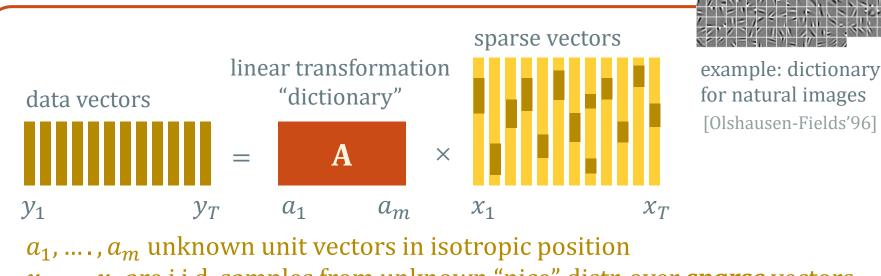
global optima of polynomial  $\langle T, x^{\otimes d} \rangle = \sum_{i=1}^{m} \langle a_i, x \rangle^d + \langle Z, x^{\otimes d} \rangle$ over unit sphere  $\approx_{\varepsilon} \{\pm a_1, \dots, \pm a_m\}$ 

also:  $\exists$  short sum-of-squares certificate for this fact

*but:* local behavior controlled by error *Z* → local search algorithms fail (also *simultaneous diagonalization* fails)

# results: dictionary learning (aka sparse coding)

*application:* machine learning (*feature extraction*) neuroscience (*model for visual cortex*)



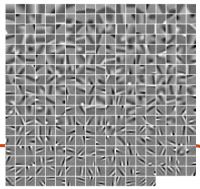
 $x_1, ..., x_t$  are i.i.d. samples from unknown "nice" distr. over *sparse* vectors (only small correlations between coord's)

*goal:* given data vectors  $y_1, \ldots, y_T$ , reconstruct A

# reduces to tensor decomposition with *spectral error* controlled by *sparsity*

[Arora-Ge-Moitra, Agarwal-Anandkumar-Jain-Netrapalli-Tandon]

*previous methods (local search):* only very sparse vectors, up to  $\sqrt{n}$  non-zeros



# results: dictionary learning (aka sparse coding)

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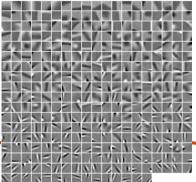


 $a_1, \ldots, a_m$  unknown unit vectors in isotropic position  $x_1, \ldots, x_t$  are i.i.d. samples from unknown "nice" distr. over **sparse** vectors (only small correlations between coord's)

*goal:* given data vectors  $y_1, \ldots, y_T$ , reconstruct A

[Arora-Ge-Moitra, Agarwal-Anandkumar-Jain-Netrapalli-Tandon]

previous methods (local search): only very sparse vectors, up to  $\sqrt{n}$  non-zeros sum-of-squares method: full sparsity range, up to constant fraction non-zeros (quasipolynomial-time for sparsity o(1); polynomial-time for  $n^{-\varepsilon}$ )



example: dictionary for natural images [Olshausen-Fields'96] simplified problem  $\langle T, x^{\otimes 3} \rangle \approx_{s}$  $a_1, \ldots, a_n \in \mathbb{R}^n$  orthonormal,  $Z \in \mathbb{R}^{n^3}$  with  $||Z||_{\text{spectral}} \leq \varepsilon$  $\max(a_i, x)$ given tensor  $T = Z + \sum_{i} a_i^{\otimes 3}$ , maximize polynomial  $\langle T, x^{\otimes 3} \rangle$  $a_2$  $a_{4}$  $a_{5}$ over unit sphere  $S^{n-1} \subseteq \mathbb{R}^n$  $a_2$ deg-k sum-of-squares algorithm computes "pseudo distribution  $D: S^{n-1} \to \mathbb{R}$ " in time  $n^{O(k)}$ behaves like *deg-k part of density of distribution* supported on solutions to  $C = \{ \langle T, x^{\otimes 3} \rangle \ge 1 - \varepsilon, \|x\|^2 = 1 \}$ i.e., D passes all tests derivable from C by deg-k SOS proof system concretely,  $\int_{S^{n-1}} D \cdot \left[ P^2 \cdot \left( \langle T, x^{\otimes 3} \rangle - (1-\varepsilon) \right) + Q^2 \right] \ge 0$  whenever deg *P*, deg  $Q \le k$ 

# want: rounding algorithm

given pseudo-distribution *D*, compute solution to constraints *C* approach: [Barak-Kelner-S.'14]

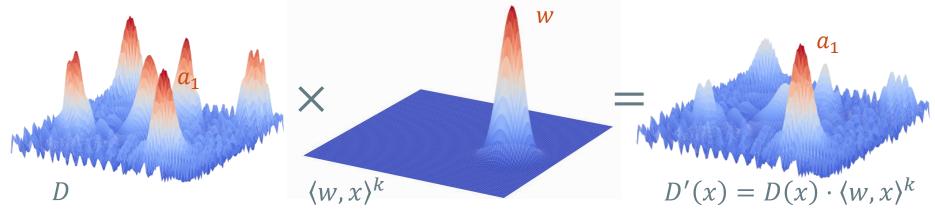
first analyze algorithm when *D* is deg-*d* part of actual distribution

simplified problem

 $a_1, ..., a_n \in \mathbb{R}^n$  orthonormal,  $Z \in \mathbb{R}^{n^3}$  with  $||Z||_{\text{spectral}} \leq \varepsilon$ given tensor  $T = Z + \sum_i a_i^{\otimes 3}$ , solve  $\mathcal{C} = \{\langle T, x^{\otimes 3} \rangle \geq 1 - \varepsilon, ||x||^2 = 1\}$ 

**assume:** D is deg-k part of density supported on solutions to  $\mathcal{C}$ 

*algorithm:* (1) reweigh *D* by  $\langle w, x \rangle^k$  for  $k \approx \log n$  and random unit vector *w* (2) output top eigenvector of resulting covariance matrix



### analysis

property of Gaussian distribution: with probability  $\geq 1/n^{O(1)}$  $\langle w, a_1 \rangle^2 \geq 2 \cdot \max_{i \geq 1} \langle w, a_i \rangle^2$ 

→ increase probability mass on  $a_1$  by factor  $2^k$  relative to other spikes → for  $k = \log n$ , almost all mass on  $a_1$  → can recover  $a_1$  from covar. matrix

#### simplified problem

 $a_1, ..., a_n \in \mathbb{R}^n$  orthonormal,  $Z \in \mathbb{R}^{n^3}$  with  $||Z||_{\text{spectral}} \leq \varepsilon$ given tensor  $T = Z + \sum_i a_i^{\otimes 3}$ , solve  $\mathcal{C} = \{\langle T, x^{\otimes 3} \rangle \geq 1 - \varepsilon, ||x||^2 = 1\}$ 

## what does it mean to efficiently certify that C has only few solutions?

derive inequality  $\sum_{i} \langle a_{i}, x \rangle^{k} \ge (1 - 2\varepsilon)^{k}$  for  $k = \log n$  from constraints C in **deg-k SOS proof system**  $(\sum_{i} y_{i}^{k})^{1/k} \approx \max_{i} y_{i}$ 

### derivation sketch

 $\begin{aligned} & \operatorname{from} \begin{cases} \|x\|^2 = 1 \\ \langle T, x^{\otimes 3} \rangle \geq 1 - \varepsilon \end{cases} & \operatorname{derive} \quad \sum_i \langle a_i, x \rangle^3 \geq 1 - 2\varepsilon \\ & \operatorname{using} \|T - \sum_i a_i^{\otimes d}\|_{\operatorname{spectral}} \leq \varepsilon \text{ (SOS captures eigenvalue bounds)} \end{aligned}$  $\begin{aligned} & \operatorname{from} \begin{cases} \|x\|^2 = 1 \\ \sum_i \langle a_i, x \rangle^3 \geq 1 - 2\varepsilon \end{cases} & \operatorname{derive} \quad \sum_i \langle a_i, x \rangle^k \geq (1 - 2\varepsilon)^k \text{ for all } k \geq d \\ & \operatorname{using} \left(\sum_i y_i^k\right) \cdot \left(\sum_i y_i^2\right)^k - \left(\sum_i y_i^3\right)^k \text{ is sum of squares,} \\ & \operatorname{choosing} y_i = \langle a_i, x \rangle, \text{ and } \operatorname{using} \sum_i \langle a_i, x \rangle^2 = \|x\|^2 \end{aligned}$ 

#### summary

polynomial optimization is easy if we can *certify* that there are *only few good solutions* 

(derive constraint of form  $\sum_i \langle a_i, x \rangle^k$  for  $k > \log \#$  solutions)

## open questions / subsequent work

sum of squares useful for other machine learning problems?
tensor prediction [Barak-Moitra]
overcomplete average-case 3-tensor decomposition [Ge-Ma]
can sum of squares lead to fast algorithms?
tensor principal component analysis [Hopkins-Shi-S.]
overcomplete average-case 3-tensor decomp. [Hopkins-Schramm-Shi-S.]
planted sparse vector [Hopkins-Schramm-Shi-S.] Thank you!